## Generational GC

#### The Generational Hypothesis

Hypothesis: most objects "die young"

- i.e., they are only needed for a short time, and can be collected soon
- A great example of empirical systems work
- Found to hold for programs in practice
  - Regardless of programming style! (functional, imperative, OO, etc.)

#### The Generational Hypothesis

```
(define (euclidian-norm v)
(sqrt (foldl + 0 (map square v))))
```

- square produces floats which are only used until the addition
- map produces a whole list which is only used until we fold over it
- + produce a whole lot of intermediate floats which are only used for the next addition

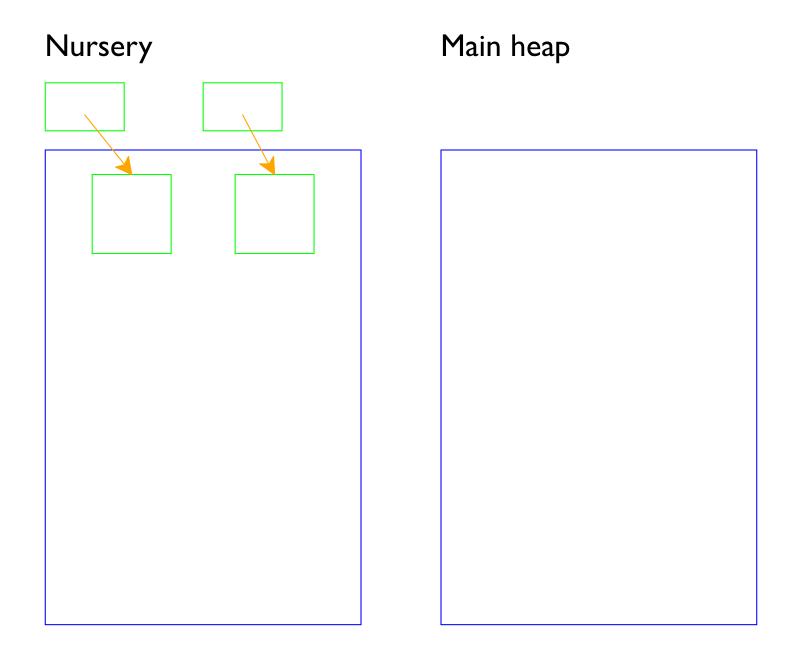
#### The Generational Hypothesis

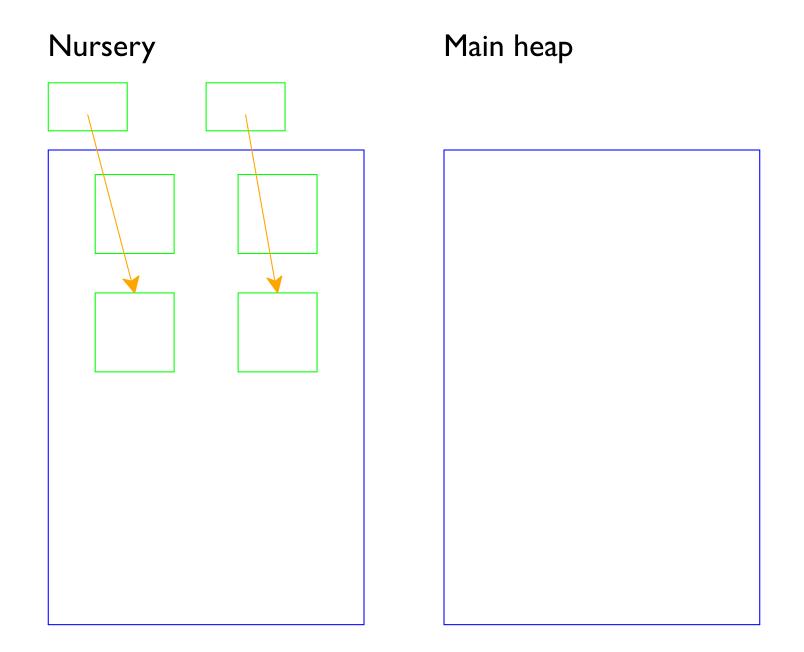
So, how can we use this to guide GC design?

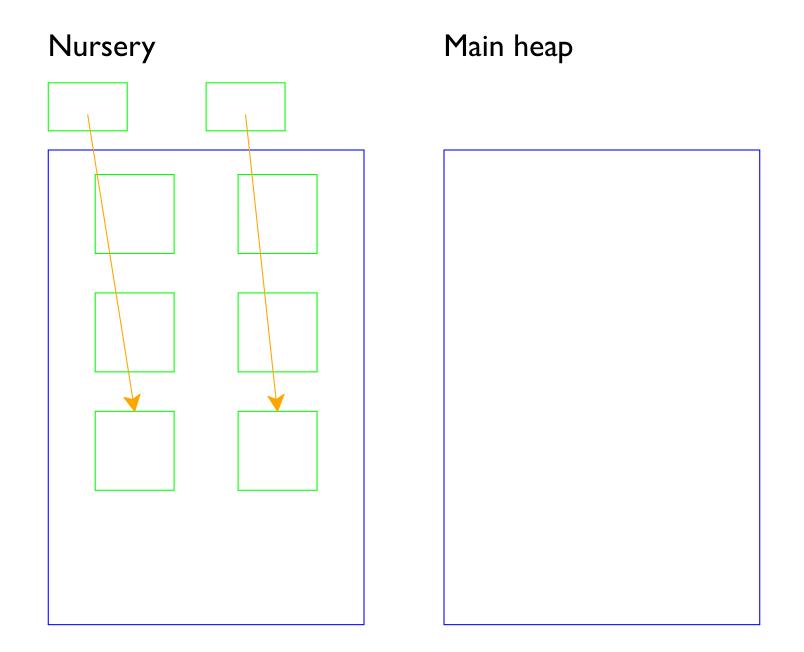
- Young objects have a high chance of being garbage
  - But they're only a small portion of all objects
  - So if we focus our efforts on them, can free a lot of space in not much time!
- Have a separate region for young objects: nursery
  - Allocate new objects there
  - When we run out of space, collect there (only!)
  - When objects get old enough, so less chance of dying, migrate them to the main heap
  - When we run out of space in the main heap, only then do we GC it

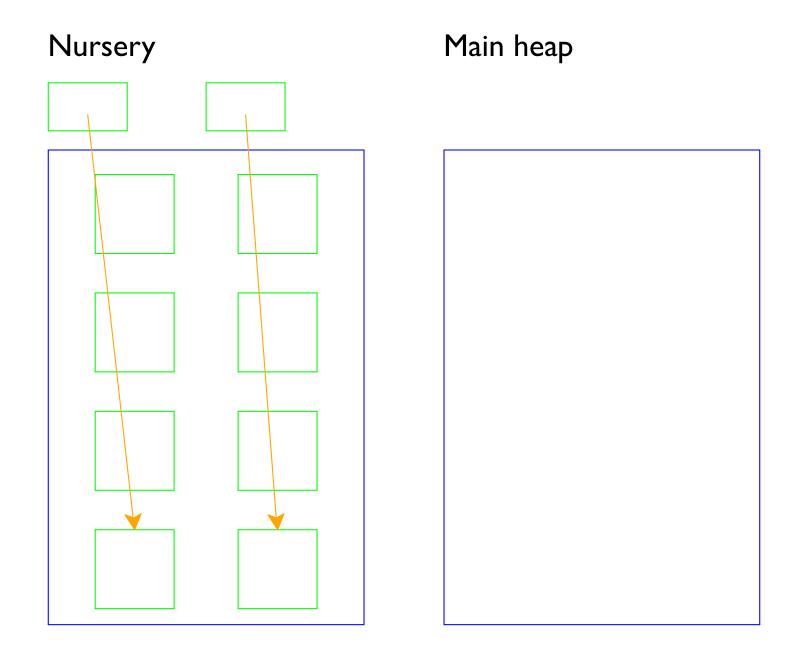
#### Ok, so how to GC?

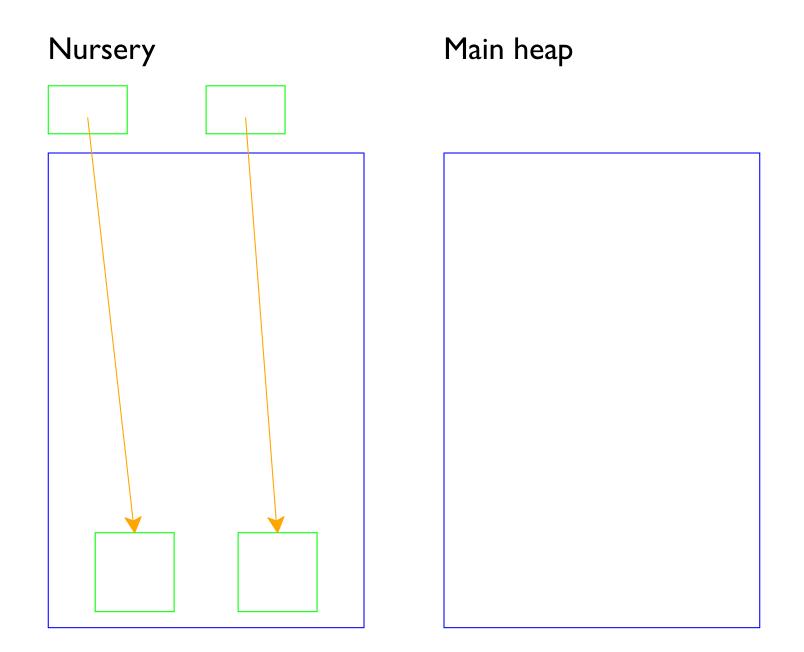
- Clearly want copying GC for nursery
  - Cost proportional to # of live objects
    - Which should be few in nursery
  - Use the main heap as to-space!
    - And don't swap spaces! Always in the same direction
    - Can use the whole heap with a copying GC!
  - Survivors get copied out of the nursery naturally!
    - Surviving a GC = old enough to move out
- What about main heap?
  - Doesn't matter as much (stay tuned)

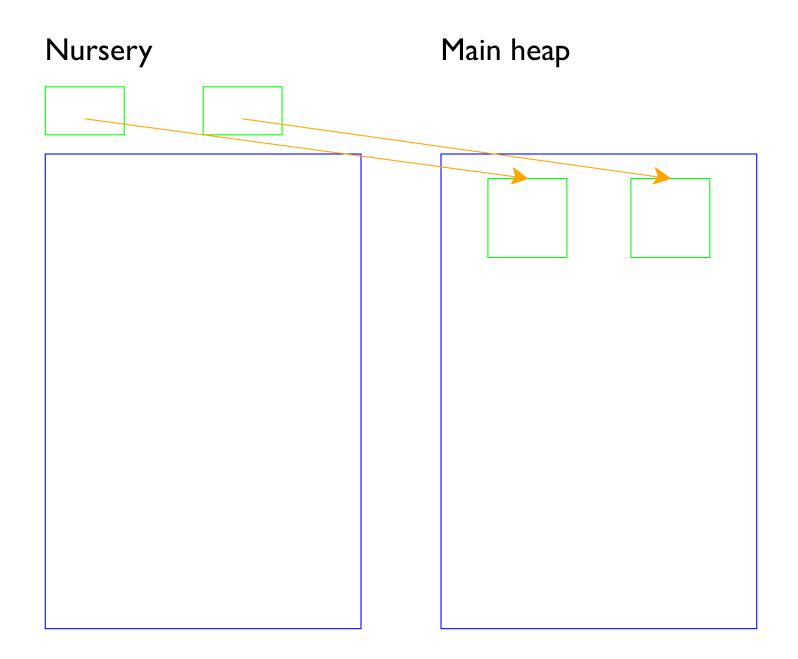


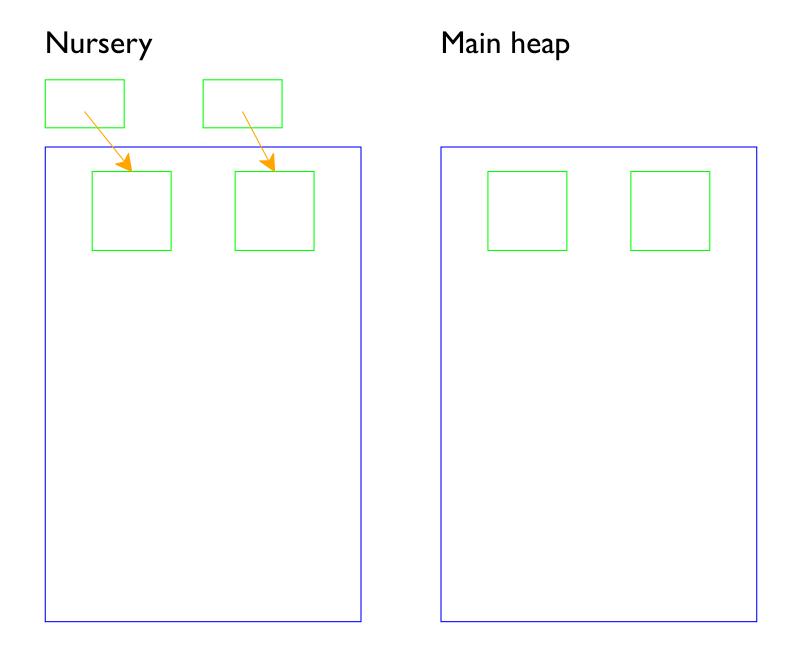


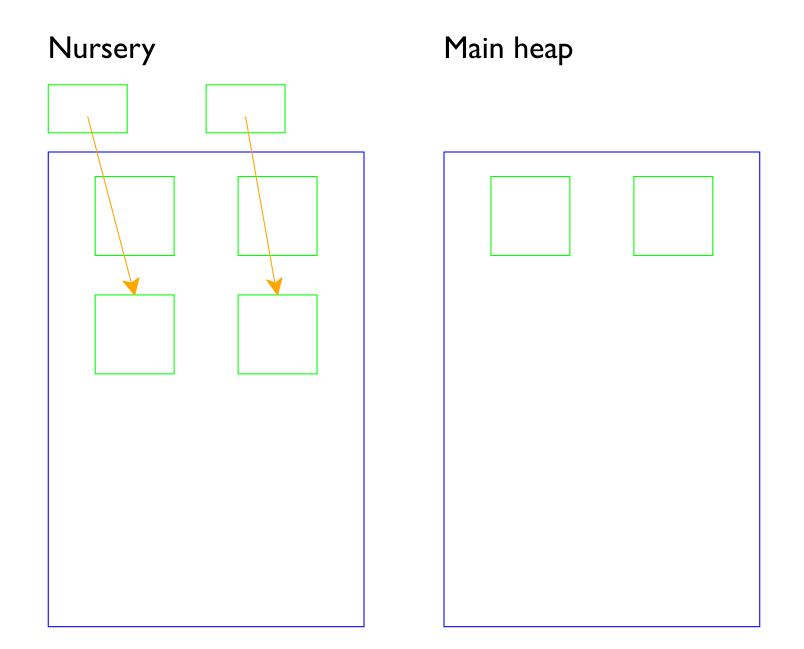


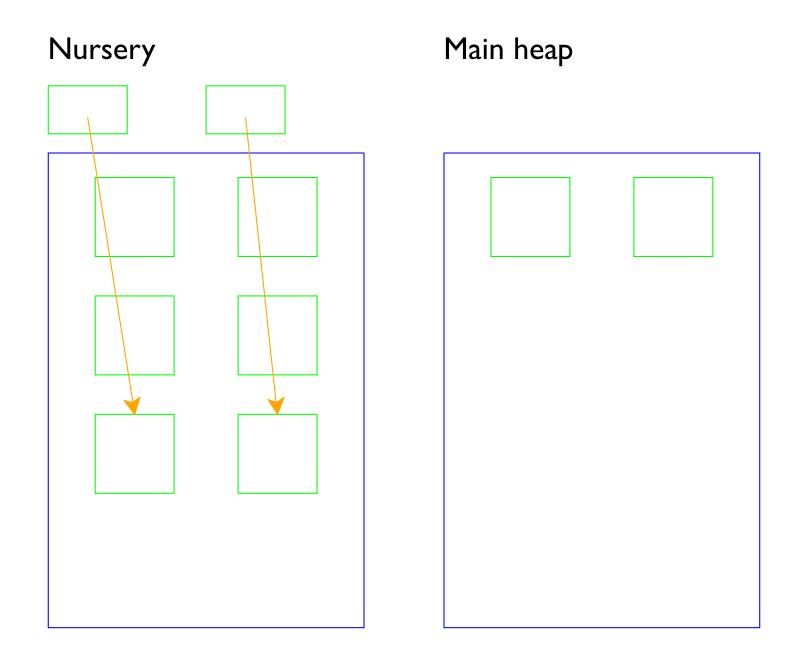


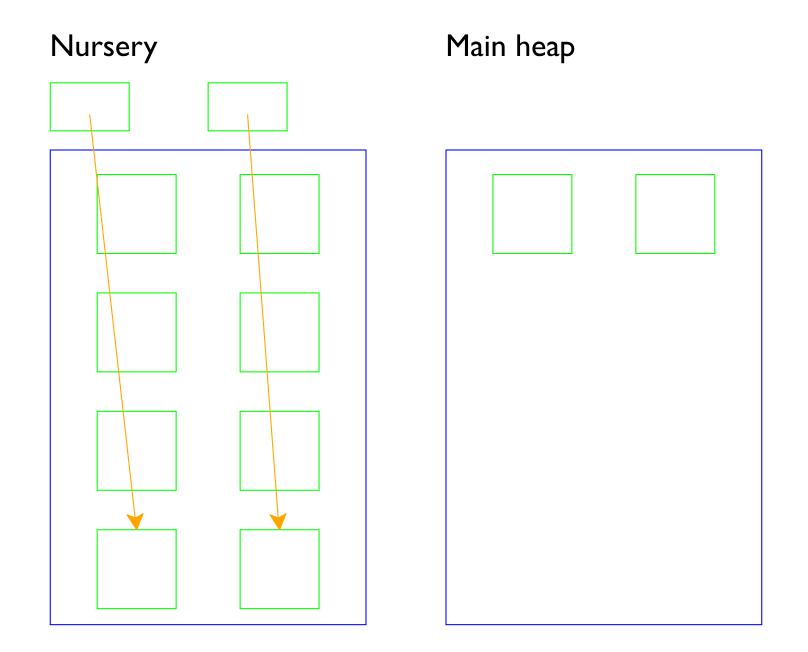


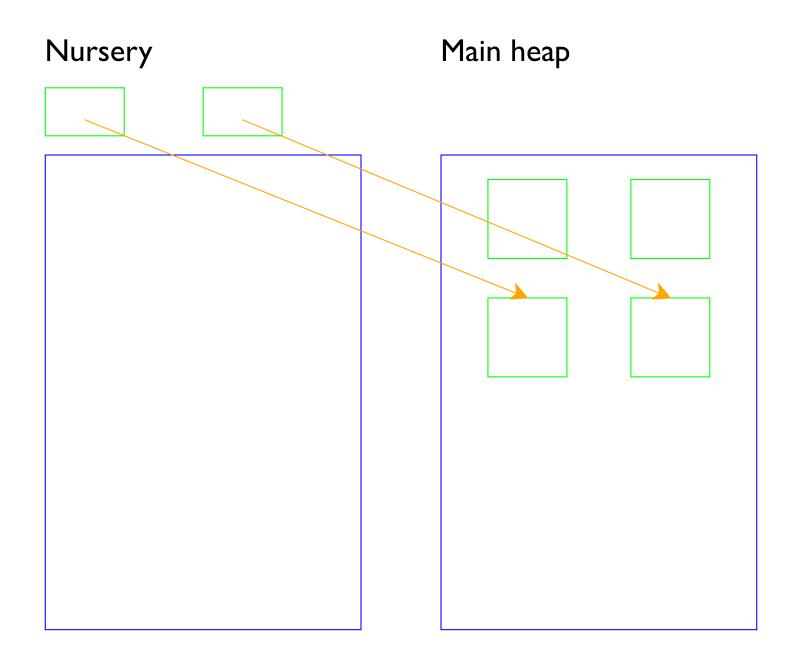


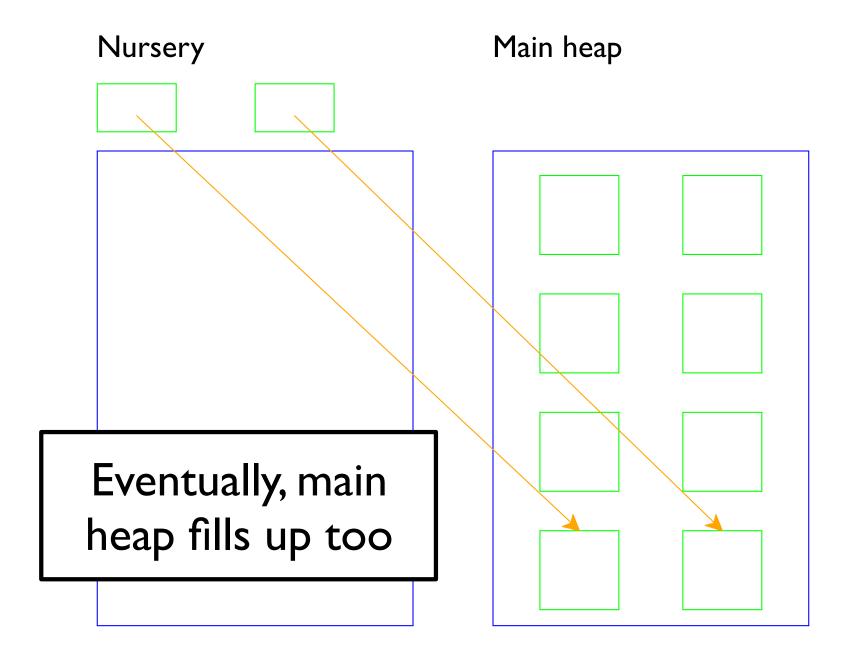


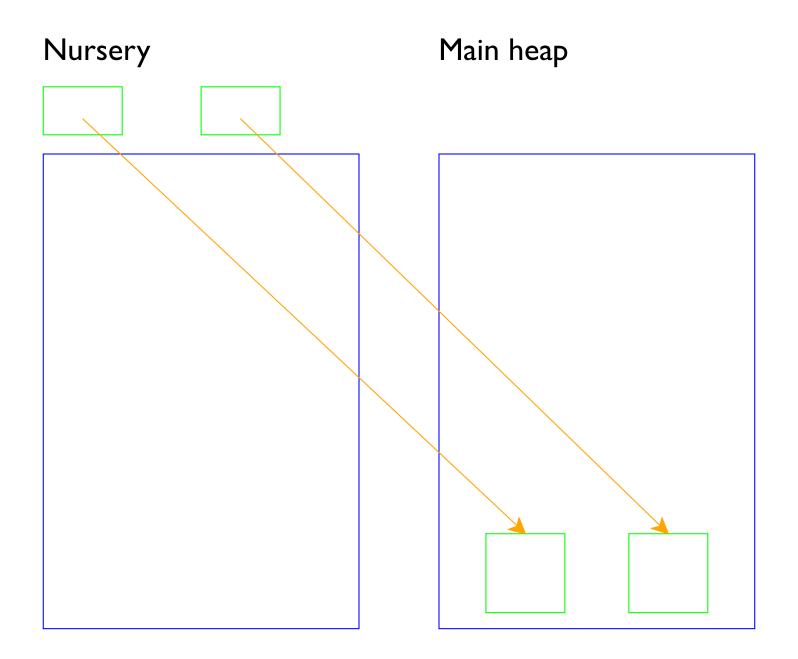


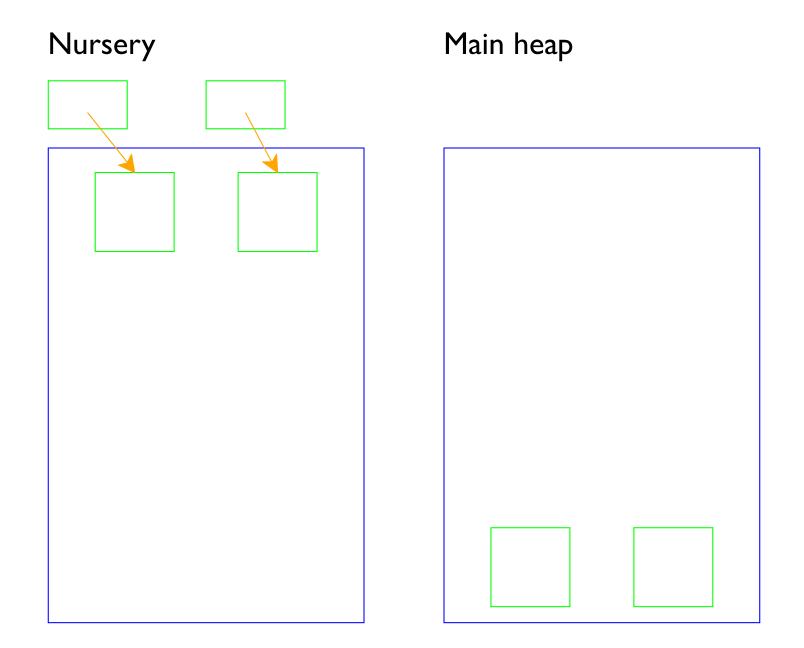


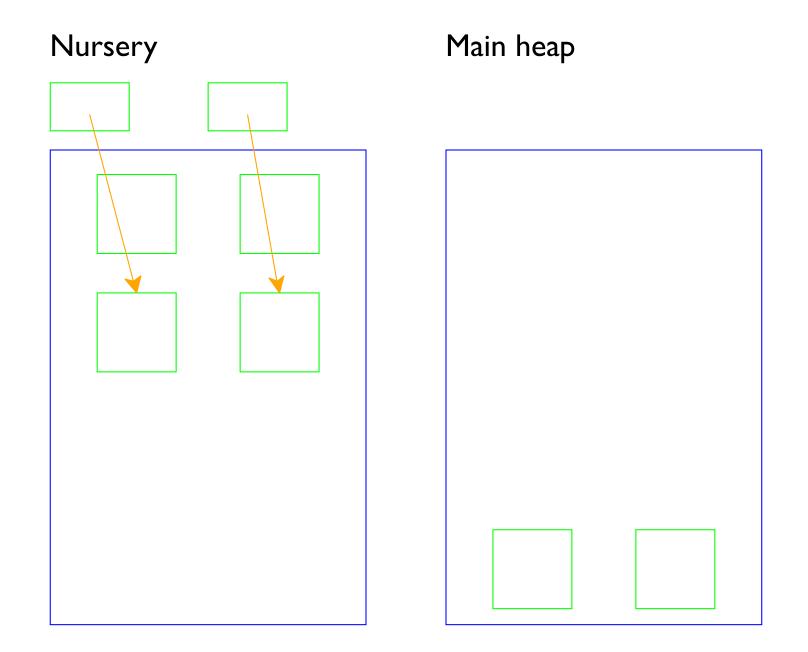


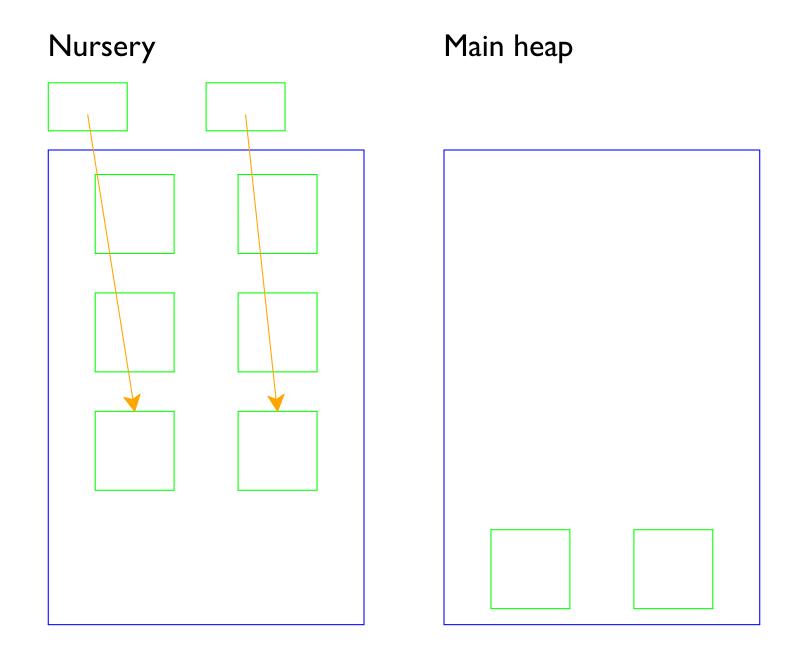


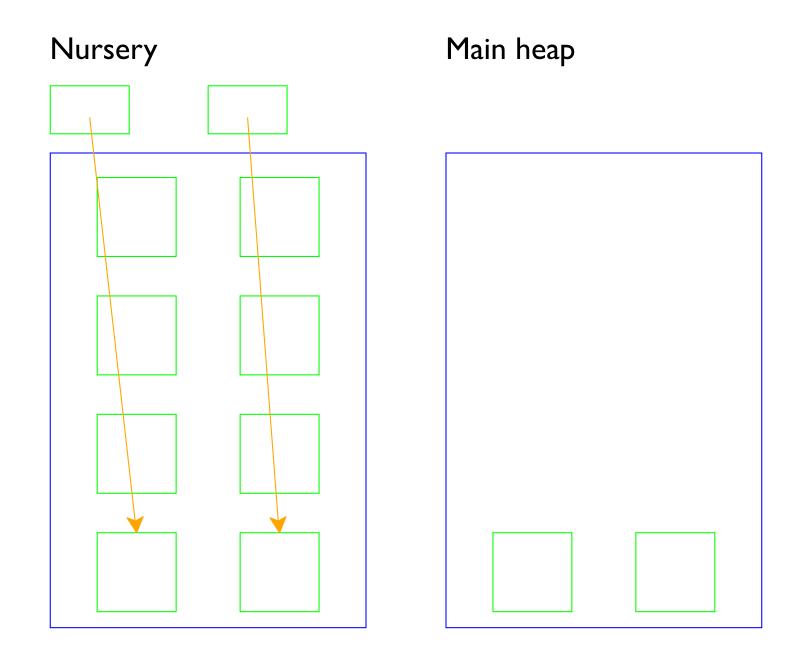


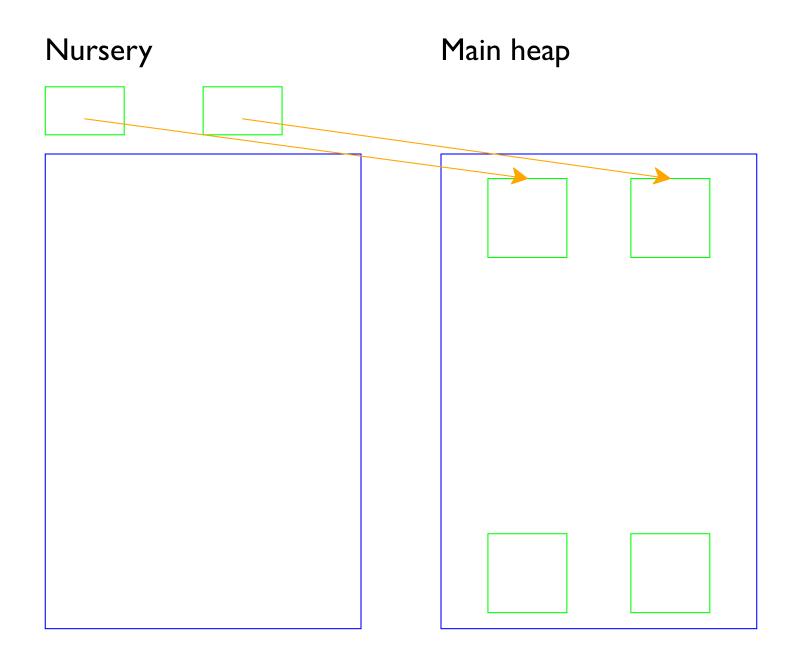












#### If the generational hypothesis holds

- Then we'll collect the nursery pretty often
- But only collect the main heap rarely

#### Collecting the nursery is cheap

- Can keep it fairly small, so not many objects
- Will be mostly dead objects, and copying GC has cost proportional to live objects!
- → can afford to do it often; pauses will be short

#### Collecting the main heap is slow

- It will be large; needs to hold all our data
- · A lot of it will be live, and will need to be marked/copied
- → but that's ok, don't do it very often

#### The Snag

- When we GC the nursery, what do we use as roots?
- Want to use registers, globals, etc. Sure.
- But we may also have pointers to nursery objects from the main heap!
  - These nursery objects may be live too!
- Only matters if the object in the main heap is live
  - But can't know unless we mark all of main heap
  - Which is what we were trying to avoid in the first place!

#### The Solution

- Track pointers from main heap to nursery
  - And just assume they're live, conservative
  - So treat them as roots
- How do we keep track? Write barriers
  - Can only get a pointer from main heap to nursery by using mutation
  - Can't have an old object point to a new one naturally
    - Latter wasn't around when the old one was allocated!
  - So whenever we mutate, we look out for that case, and keep track
- Doesn't happen much in practice
  - Requires a particular mutation pattern
  - So ok to be conservative

#### **Variants**

- Can have N generations
  - Containing progressively older and older objects
- Can migrate generations only after surviving N GCs
  - Reduces the number of short-lived objects that only survive because we happened to GC during their (short) lifetime
  - Can use some spare bits in the objects to store count (bit overloading, think mark bit)

#### In Practice

- The vast majority of production GCs are generational in some form
  - It's that good
- Generational hypothesis: self-fulfilling prophecy / virtuous circle
  - Generational GC is efficient because most objects are short-lived
  - Generational GC makes short-lived objects cheap
  - Programmers use more short-lived objects because they're cheap
  - Lather, rinse, repeat

#### History

- Came from the Self language (late 80s, early 90s)
- Self is little-known today, but hugely influential
  - JS is basically Self
  - Implementation technology (JITs, PICs, OSR, adaptive deopts, etc.) is used all over the place

#### Charlie on the MTA

- 1940s: Massachusetts Transit Association (MTA) has both entry and exit fares on trains
- 1949: Walter A. O'Brien runs for mayor of Boston
  - His campaign: get folk singers to write songs about items on his platform
  - Then blare them from a truck going around town (got fined 10\$ for that)
- One such song is about Charlie, who has enough money to get on the train, but not off
  - So he never returned
- 1959: The song itself becomes a hit
- 2006: MBTA (former MTA) introduces the Charlie card
- 2009: I move to Boston, and finally get the Charlie/Cheney joke

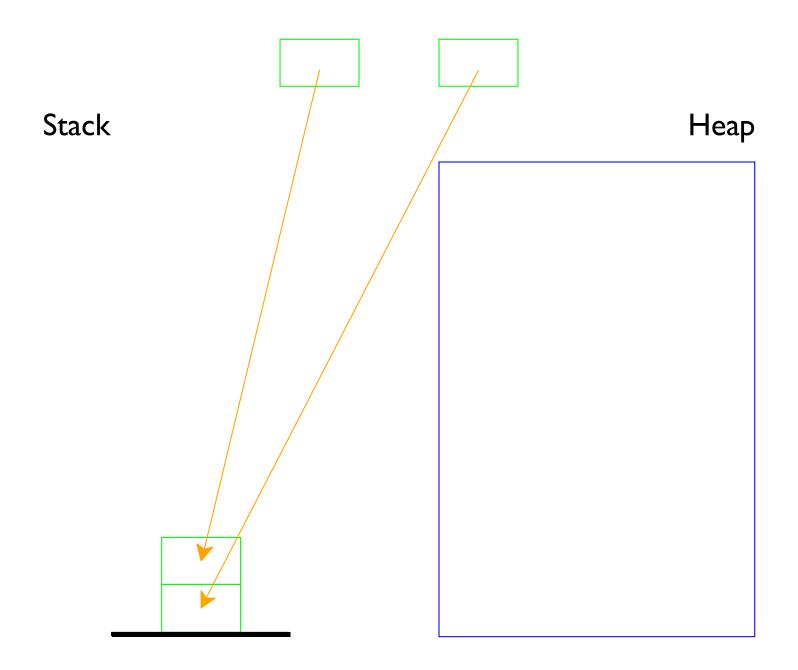
#### Continuation-Passing Style

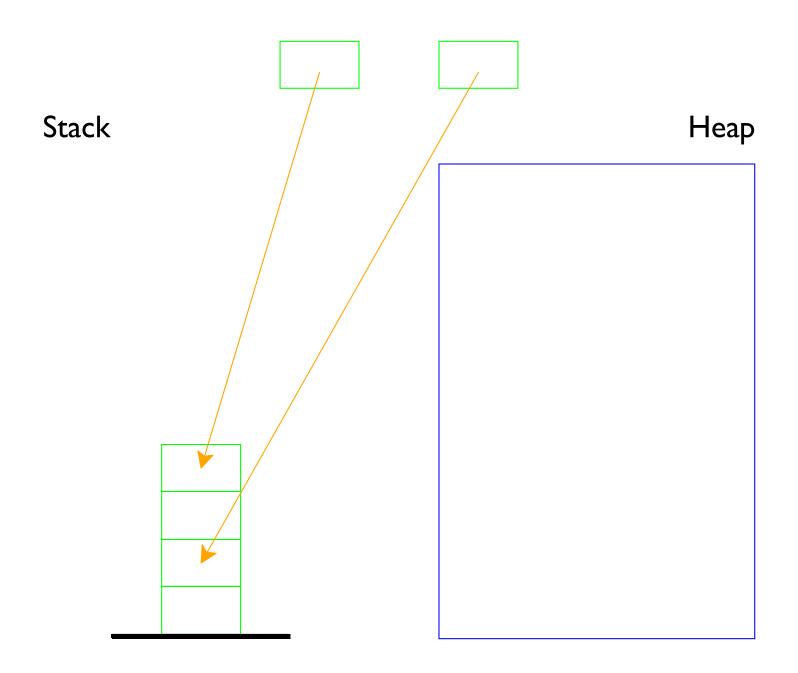
- We've seen continuation-passing style
  - Wrote an interpreter that way
- In CPS, the last thing a function does is always call another function
  - If the function is done with its own work, it calls its continuation
  - Otherwise, passes it along to whoever it calls
  - So our functions never return
  - They just keep calling until the end, then everything just returns all at once
- Compilers for language with higher-order functions often convert object programs to CPS inside the compiler
  - Easier to implement control (e.g., return, exceptions, etc.)
  - Makes a lot of transformations and optimizations easier

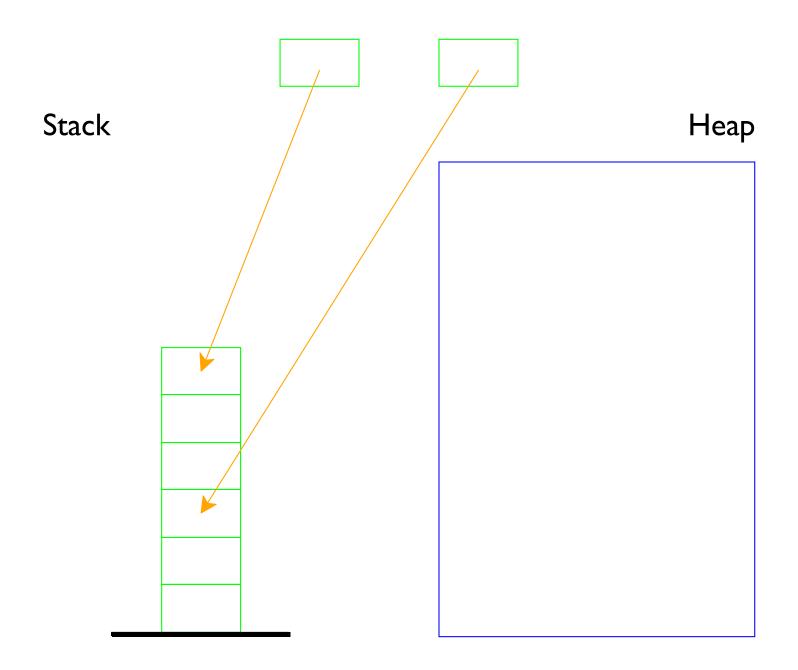
- Two-space copying collection = Cheney's algorithm
- Our object programs are in CPS
  - Their functions never return!
  - Just like Charlie!
- So their stack just keeps growing and growing
  - (In CPS, all calls are tail calls, so could reuse stack frames and solve that problem, which is how most compilers solve the problem. But not here!)

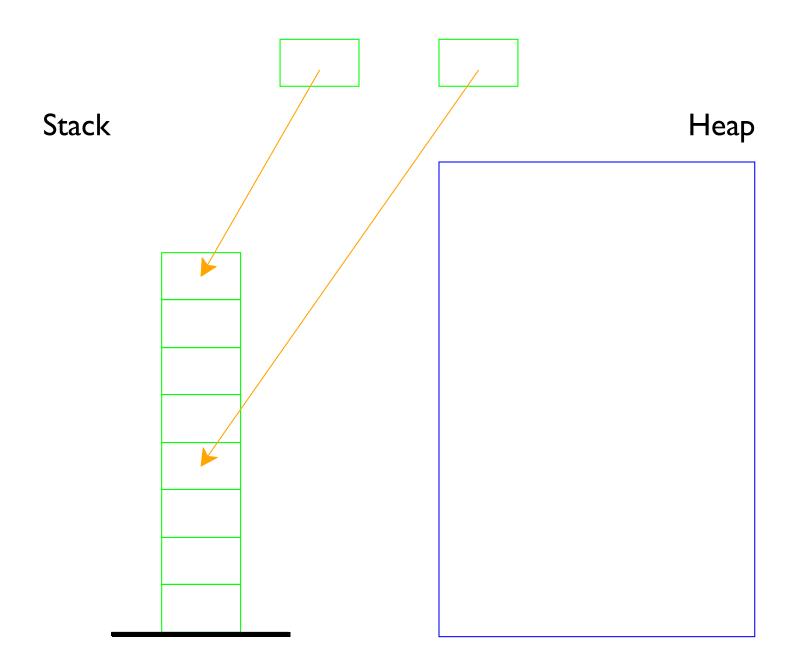
# Key idea: use the stack as the nursery for a generational GC!

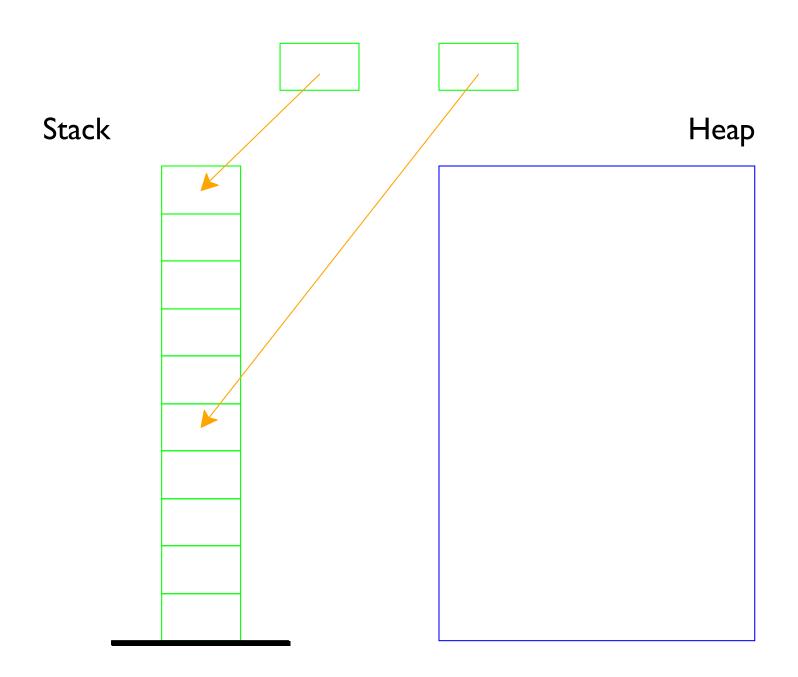
- Grow the stack until we run out of stack space
- Then GC, copying live objects into the heap
- Then restart the stack from 0!
  - Ohthing ever returns, so we don't need return addresses!
  - And all the useful data has been copied away!
- · Analogy: instead hopping on a trampoline on every function call
  - We sometimes jump over the Empire State Building
  - (Trampolines are a canonical solution for "reusing" stack frames for tail calls)

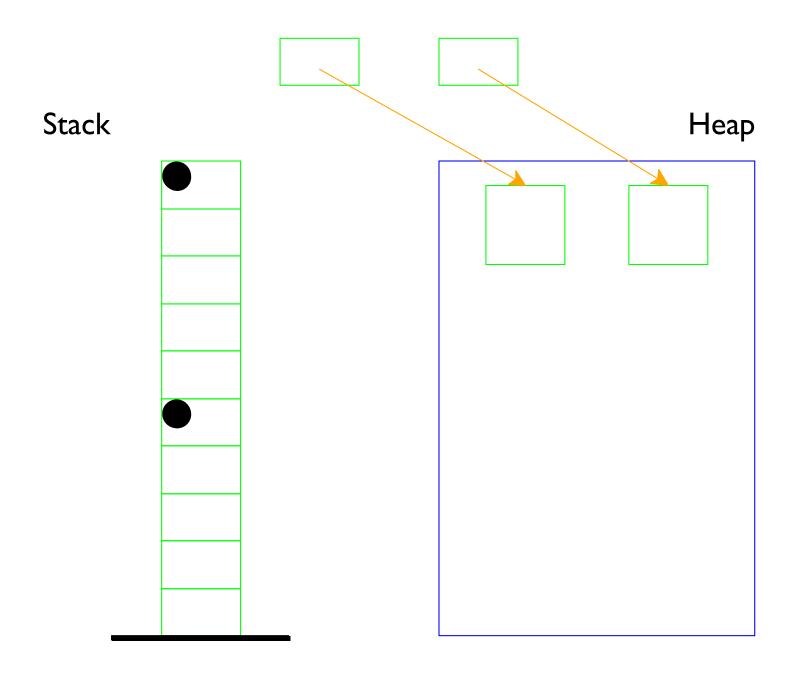


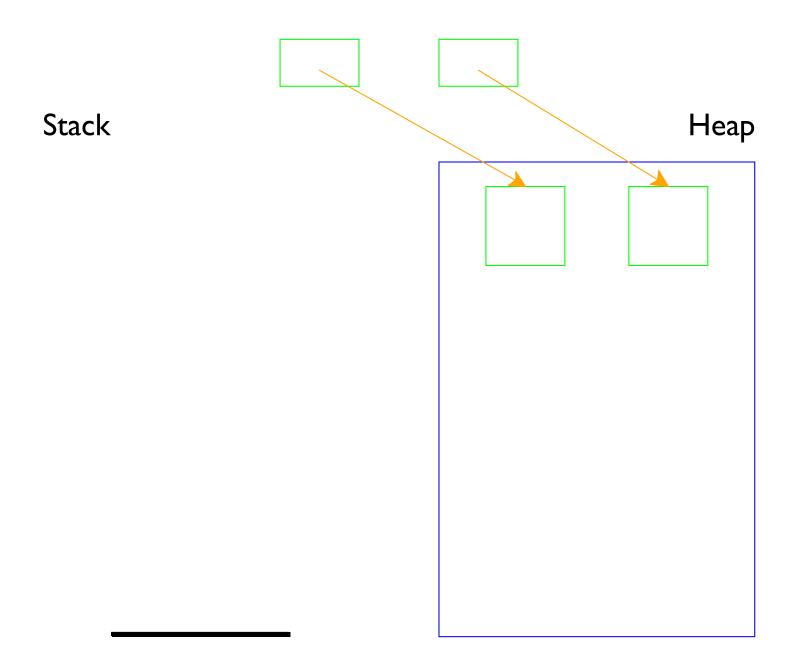


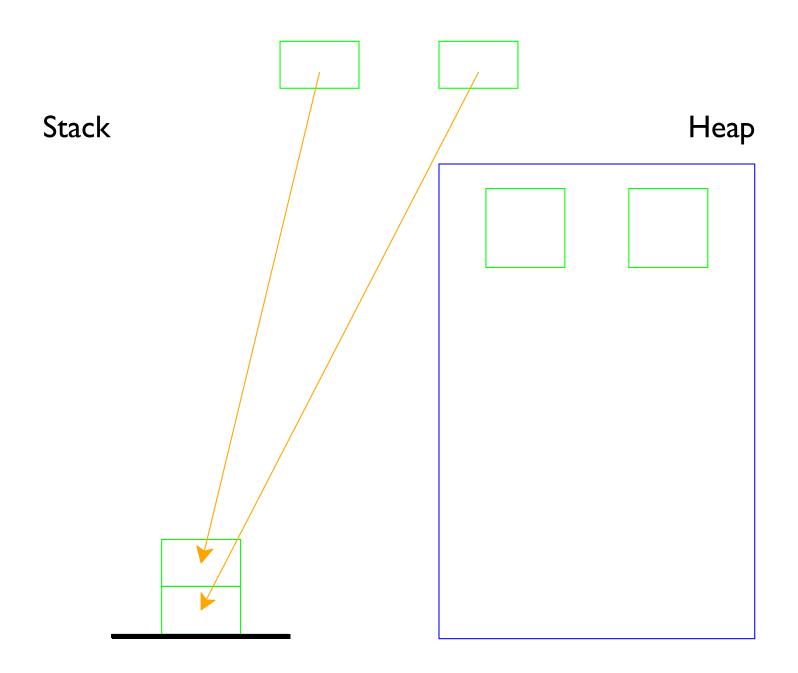


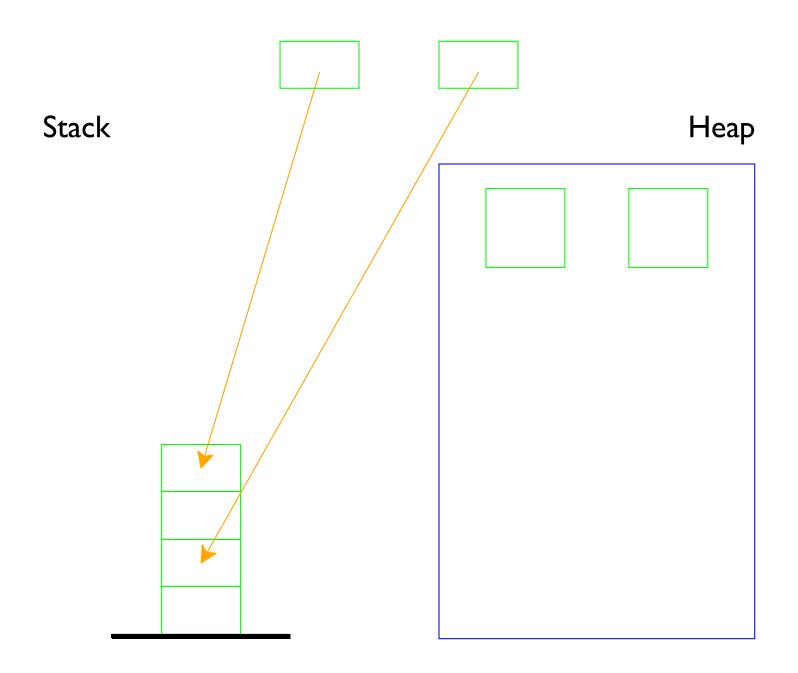


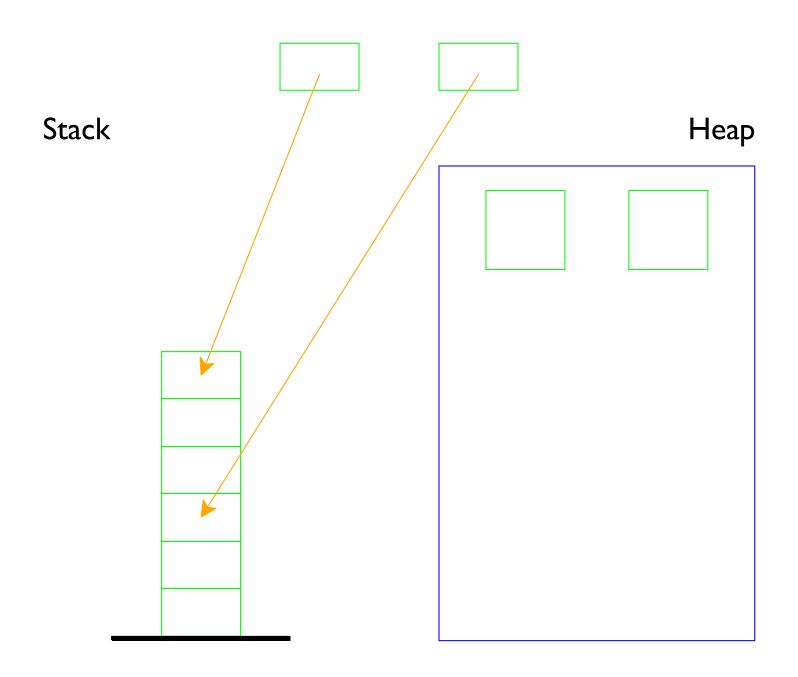


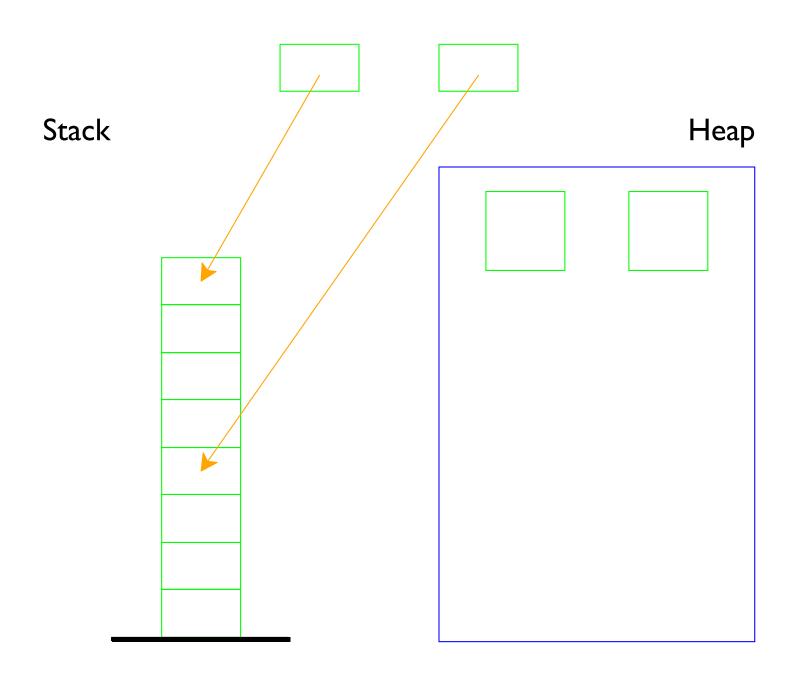


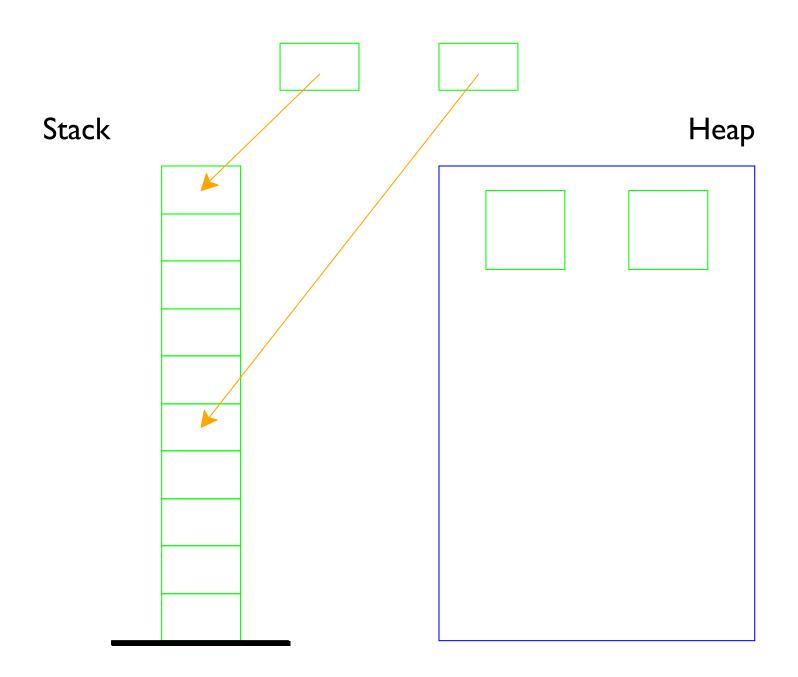


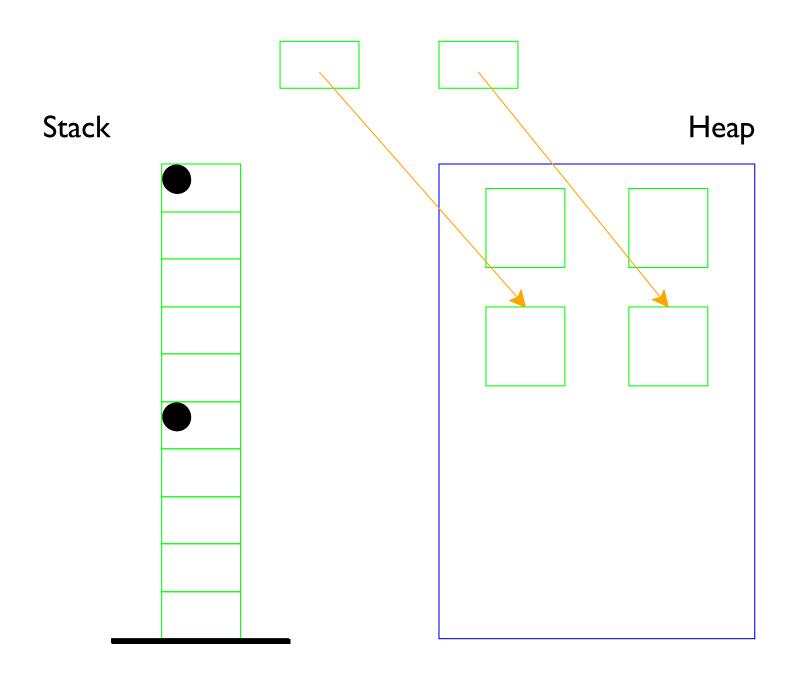


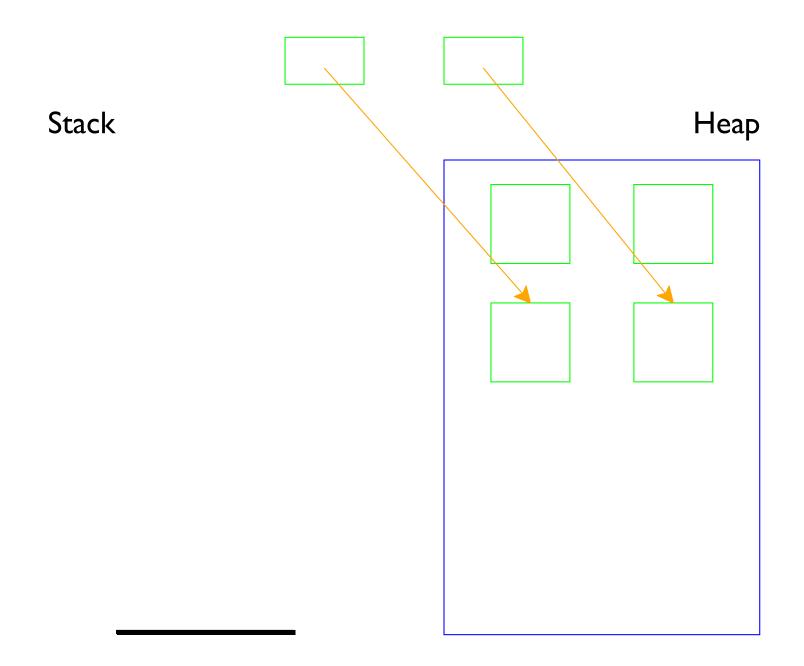












### For More Information

- Chicken Scheme uses this implementation strategy
- CONS Should Not CONS Its Arguments, Part II: Cheney on the M.T.A., Henry Baker, 1994
  - O http://home.pipeline.com/~hbaker1/CheneyMTA.html

