# Typed Imperative Programming 

## CS 211

Winter 2020

## Initial code setup

The code in this course is available online. To download a copy of this lecture into your Unix shell account:
\% cd cs211
\% curl -k \$URL211/lec/02typed_imp.tgz | tar zxv
\% cd 02typed_imp

## Initial code setup

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- The curl(1) command downloads a URL and prints its contents to its standard output (stdout).


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- The curl(1) command downloads a URL and prints its contents to its standard output (stdout).
- The $\operatorname{tar}(1)$ command extracts various forms of archives. (The "(1)" means you can get help by running man 1 tar.)


## Initial code setup

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\% cd cs211
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- The curl(1) command downloads a URL and prints its contents to its standard output (stdout).
- The tar(1) command extracts various forms of archives. (The "(1)" means you can get help by running man 1 tar.)
- The | character is a Unix "pipe," which attaches the first command's stdout to the second command's stdin.

Problem: Compute the $n$th Fibonacci number

## Definition

$$
f i b(n) \begin{cases}n & \text { if } n<2 \\ \text { fib }(n-2)+f i b(n-1) & \text { otherwise }\end{cases}
$$

## Definition

$$
f i b(n) \begin{cases}n & \text { if } n<2 \\ \operatorname{fib}(n-2)+f i b(n-1) & \text { otherwise }\end{cases}
$$

| $n$ | fib $(n)$ |
| :---: | ---: |
| 0 | 0 |
| 1 | 1 |
| 2 | 1 |
| 3 | 2 |
| 4 | 3 |
| 5 | 5 |
| 6 | 8 |
| 7 | 13 |
| 8 | 21 |

## In C (don't do this at home!)

long fib(int n)
\{

$$
\text { return } n<2
$$

? n
: fib(n - 2) + fib(n - 1);
\}

## In C (don't do this at home!)

long fib(int n)
\{

$$
\text { return } n<2
$$

? n

$$
: f i b(n-2)+f i b(n-1) ;
$$

\}
long fib(int n)
\{

$$
\text { return }(n<2) ? n:(f i b(n-2)+f i b(n-1)) ;
$$

\}

## In C (don't do this at home!)

long fib(int n)
\{

$$
\text { return } n<2
$$

? n
: fib(n - 2) + fib(n - 1);
\}
long fib(int n) $\{$ return $n<2 ? n: f i b(n-2)+f i b(n-1) ;\}$

## In C (don't do this at home!)

long fib(int n)
\{

$$
\text { return } n<2
$$

$$
? \mathrm{n}
$$

$$
: f i b(n-2)+f i b(n-1) ;
$$

\}
long fib(int n) \{
return $n<2$ ?n:fib
$(n-2)+f i b(n-1) ;\}$

## In C (don't do this at home!)

```
long fib(int n)
{
    return n < 2
    ? n
    : fib(n - 2) + fib(n - 1);
```

\}

Things to notice:

- Static types int ${ }^{a}$ and long $^{b}$ must be given for variables (argument n) and function results.
- This function does computation but not input/output.

[^0]
## In C (less weird but still slow [and weird]) <br> long fib(int n) <br> \{ <br> if ( n < 2) \{ return n; <br> \} else \{ <br> long $b=f i b(n-1) ;$ return a + fib(n - 1); <br> \} <br> \}

## In C（less weird but still slow［and weird］）

long fib（int n）
\｛

$$
\begin{aligned}
& \text { if }(n<2)\{ \\
& \quad \text { return } n ; \\
& \} \text { else }\{ \\
& \quad \text { long } b=f i b(n-1) ; \\
& \text { return } a+f i b(n-1) ;
\end{aligned}
$$

\}
Syntax of if：


```
In C (less weird but still slow [and weird])
long fib(int n)
{
    if (n<2) {
        return n;
} else {
    long b = fib(n - 1);
    return a + fib(n - 1);
}
\}
```

Syntax of variable definition:

$$
\langle\text { type }\rangle\langle\text { var-name }\rangle=\langle\text { init-expr }\rangle ;
$$

Semantics: allocate space named $\langle$ var-name $\rangle$ for a value of type $\langle$ type $\rangle$; evaluate $\langle$ init-expr $\rangle$ and store its result there.

```
In C (less weird but still slow [and weird])
long fib(int n)
{
    if (n< 2) {
        return n;
} else {
    long b = fib(n - 1);
    return a + fib(n - 1);
}
\}
```

Syntax of return:

```
return \langleresult-expr\rangle;
```

Semantics: evaluate $\langle$ result-expr $\rangle$, then return that value from this function immediately.

## In C (less weird but still slow [and weird])

long fib(int n)
\{

$$
\begin{aligned}
& \text { if }(n<2)\{ \\
& \text { return } n \text {; } \\
& \text { \} else \{ } \\
& \quad \text { long b = fib }(n-1) ; \\
& \text { return } a+\operatorname{fib}(n-1) ;
\end{aligned}
$$

\}
More syntax of if:

$$
\begin{aligned}
& \text { if } \begin{array}{l}
(\langle\text { test-expr }\rangle) \\
\text { \} \{then-stms }\rangle
\end{array}
\end{aligned}
$$

## Everything nests



## Everything nests

```
if (\langlefirst-test-expr\rangle) { // But don't write this.
    A-stms\rangle
} else {
    if (\langlesecond-test-expr\rangle) {
        <B-stms\rangle
    } else {
        <C-stms\rangle
    }
}
if (\langlefirst-test-expr\rangle) { // Do write this.
    |A-stms\rangle
} else if (\langlesecond-test-expr\rangle) {
    <B-stms\rangle
} else {
    <C-stms\rangle
}
```

Problem: It's super slow

## Solution: Mutation (and iteration)

## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 ; \\
& \text { int } b=8 ; \\
& \text { int } c ; \\
& c=a+b ; \\
& a=b ; \\
& b=c ; \\
& c=a+b ; \\
& a=b ; \\
& b=c ; \\
& c=a+b ; \\
& a=b ; \\
& b=c ;
\end{aligned}
$$

## The essence of imperative programming

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& b=c ; \\
& c=a+b ; \\
& a=b ; \\
& b=c ;
\end{aligned}
$$

## The essence of imperative programming

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\begin{aligned}
& \text { int } a=5 ; \\
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& \text { int } c ;
\end{aligned}
$$

$$
\begin{aligned}
& c=a+b ; \\
& a=b ; \\
& b=c ;
\end{aligned}
$$

$$
c=a+b ;
$$

$$
a=b ;
$$

$$
b=c ;
$$

$$
c=a+b ;
$$

$$
a=b ;
$$

$$
b=c ;
$$

ax00000005

- $0 \times 00000008$


## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

a $0 \times 00000005$
b $0 \times 00000008$
c

- $c=a+b ;$
a = b;
b = c;

$$
\begin{aligned}
& c=a+b ; \\
& a=b ; \\
& b=c ;
\end{aligned}
$$

$$
c=a+b ;
$$

$$
a=b ;
$$

$$
b=c ;
$$

## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

ax00000005

- $0 \times 00000008$
c $0 \times 0000000 \mathrm{D}$
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;


## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

a $0 \times 00000008$

- 0x00000008
c $0 \times 0000000 \mathrm{D}$
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;


## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

a $0 \times 00000008$

- $0 \times 0000000 \mathrm{D}$
c $0 x 0000000 \mathrm{D}$
$c=a+b ;$
a = b;
b = c;
- $\mathrm{c}=\mathrm{a}+\mathrm{b}$;
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;


## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

a $0 \times 00000008$

- $0 \times 0000000 \mathrm{D}$
c $0 \times 00000015$
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;


## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

a 0 x0000000D
b 0x0000000D
c $0 \times 00000015$
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;

## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

ax0000000D

$$
\text { b } 0 \times 00000015
$$

$$
\text { c } 0 \times 00000015
$$

$c=a+b ;$
a = b;
b = c;

$$
\begin{aligned}
& c=a+b ; \\
& a=b ; \\
& b=c ;
\end{aligned}
$$

- $c=a+b ;$
$\mathrm{a}=\mathrm{b}$;
b = c;


## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

ax0000000D

$$
\text { b } 0 \times 00000015
$$

$$
\text { c } 0 \times 00000022
$$

$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
$\mathrm{a}=\mathrm{b}$;
b = c;

## The essence of imperative programming

$$
\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

a $0 \times 00000015$

- $0 \times 00000015$
c $0 \times 00000022$
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
$\mathrm{a}=\mathrm{b}$;
b = c;


## The essence of imperative programming

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\begin{aligned}
& \text { int } a=5 \text {; } \\
& \text { int } b=8 \text {; } \\
& \text { int }
\end{aligned}
$$

ax00000015

- $0 \times 00000022$
c $0 \times 00000022$
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;
$c=a+b ;$
a = b;
b = c;


# What's happening with variables, definitions, and assignments? 

$$
\text { int } z=5 ; \text { double } d=5 ; \text { char } c=5 ;
$$

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$$
\text { int } z=5 ; \text { double } d=5 ; \text { char } c=5 ;
$$

- Numbers (e.g., 0x000000A6 and 3.57e-99) are values.


## What's happening with variables, definitions, and assignments?

```
int z = 5; double d = 5; char c = 5;
```

- Numbers (e.g., 0x000000A6 and 3.57e-99) are values.
- An object is a location where you can store a particular type of value:
z 0x00000005
d 5.00000000 E 0000000
c $0 \times 05$
(holds an int)
(holds a double)
(holds a char)


## What's happening with variables, definitions, and assignments?

```
int z = 5; double d = 5; char c = 5;
```

- Numbers (e.g., 0x000000A6 and 3.57e-99) are values.
- An object is a location where you can store a particular type of value:
z $0 \times 00000005$
(holds an int)
d 5.00000000 E 0000000
(holds a double)
c $0 \times 05$
(holds a char)
- A variable is the name of an object (like $z$ and $d$ ).


## What's happening with variables, definitions, and assignments?

$$
\begin{aligned}
& \text { int } z=5 ; \text { double } d=5 \text {; char } c=5 \text {; } \\
& z+=c ;
\end{aligned}
$$

- Numbers (e.g., 0x000000A6 and 3.57e-99) are values.
- An object is a location where you can store a particular type of value:
z $0 \times 00000005$
(holds an int)
d 5.00000000 E 0000000
(holds a double)
c $0 \times 05$
(holds a char)
- A variable is the name of an object (like $z$ and $d$ ).
- An assignment modifies the value stored in an object.


## What's happening with variables, definitions, and assignments?

$$
\begin{aligned}
& \text { int } z=5 ; \text { double } d=5 ; \text { char } c=5 ; \\
& z+=c ;
\end{aligned}
$$

- Numbers (e.g., 0x000000A6 and 3.57e-99) are values.
- An object is a location where you can store a particular type of value:
z $0 \times 0000000 \mathrm{~A}$
d 5.00000000 E 0000000
c $0 \times 05$
(holds an int)
(holds a double)
(holds a char)
- A variable is the name of an object (like $z$ and $d$ ).
- An assignment modifies the value stored in an object.


## The other ingredient: iteration with while

Syntax:

```
while ( \(\langle\) test-expr \(\rangle\) ) \{
    〈body-stms〉
\}
```


## The other ingredient: iteration with while

Syntax:

```
while (\langletest-expr\rangle) {
    <body-stms\rangle
}
```

Semantics:

1. Evaluate $\langle$ test-expr $\rangle$ to a bool.
2. If the bool is false then the loop is finished, so jump to the next statement after the loop. after it
3. Execute $\langle$ body-stms $\rangle$.
4. Go back to step 1.

## In C, iteratively

long fib(int n)
\{

$$
\begin{aligned}
& \text { long curr = 0; } \\
& \text { long next }=1 ; \\
& \text { while (n > 0) \{ } \\
& \text { long prev = curr; } \\
& \text { curr = next; } \\
& \begin{array}{l}
\text { next += prev; } \\
\text { n } \quad=1 ;
\end{array}
\end{aligned}
$$

return curr;
\}

## In C, iteratively

long fib(int n)
\{

$$
\begin{aligned}
& \text { long curr = 0; } \\
& \text { long next = 1; } \\
& \text { while (n > 0) \{ } \\
& \text { long prev = curr; // variable definition } \\
& \text { curr = next; // assignment } \\
& \text { next += prev; // add-to } \\
& \text { n -= 1; } \\
& \text { \} }
\end{aligned}
$$

return curr;
\}

## Counting upwards

```
long fib(int n)
{
    long curr = 0;
    long next = 1;
    int i = 0;
    while (i < n) {
    long prev = curr;
    curr = next;
    next += prev;
    ++i; // equivalent to i += 1;
}
return curr;
}
```


## Counting upwards with for

```
long fib(int n)
{
    long curr = 0;
    long next = 1;
    int i = 0;
    for (; i < n; ) {
        long prev = curr;
        curr = next;
        next += prev;
    ++i;
}
```

return curr;
\}

## Counting upwards with for

```
long fib(int n)
{
    long curr = 0;
    long next = 1;
    int i = 0;
    for (; i < n; ++i) {
    long prev = curr;
    curr = next;
    next += prev;
    // ++i
}
```

return curr;

## Counting upwards with for

```
long fib(int n)
{
    long curr = 0;
    long next = 1;
    // int i = 0;
    for (int i = 0; i < n; ++i) {
    long prev = curr;
    curr = next;
    next += prev;
    // ++i
}
return curr;

\section*{Counting upwards with for}
```

long fib(int n)
{
long curr = 0;
long next = 1;
for (int i = 0; i < n; ++i) {
long prev = curr;
curr = next;
next += prev;
}
return curr;
}

```

Adding I/O

\section*{Reading user input}
\#include <stdio.h>
int main()

\section*{src/get_input.c}
\{
int \(x=0, y=0 ;\)

scanf("\%d\%d", \&x, \&y);
printf( \({ }^{\circ} \%_{\mathrm{L}} *_{\mathrm{L}} \%_{\mathrm{L}}==_{\mathrm{L}} \% \mathrm{~d} \backslash \mathrm{n}\) ", \(\mathrm{x}, \mathrm{y}, \mathrm{x}\) * y);
\}

\section*{Reading user input}
\#include <stdio.h>
int main()

> src/get_input.c
\{
\[
\text { int } x=0, y=0 ;
\]
printf("Enter \({ }^{\text {two }}\) பintegers: \({ }^{\text {匕 }}\) ");
scanf("\%d\%d", \&x, \&y);

\}
- \(\operatorname{scanf(3)}\) takes a template specifying what types of values to read and how many.

\section*{Reading user input}
\#include <stdio.h>
int main()

> src/get_input.c
\{
\[
\text { int } x=0, y=0 ;
\]

scanf("\%d\%d", \&x, \&y);
printf( \({ }^{\circ} \%_{\mathrm{L}} *_{\mathrm{L}} \%_{\mathrm{L}}==_{\mathrm{L}} \% \mathrm{~d} \backslash \mathrm{n}\) ", \(\mathrm{x}, \mathrm{y}, \mathrm{x}\) * y\()\);
- \(\operatorname{scanf(3)}\) takes a template specifying what types of values to read and how many.
- printf(3) takes a template with holes to fill in with the values of its remaining arguments.

\section*{Reading user input}
```

\#include <stdio.h>

```
int main()
src/get_input.c
\{
\[
\text { int } x=0, y=0 ;
\]
\[
\text { printf("Enter }{ }_{\text {L }} \text { two }{ }_{4} \text { integers: }{ }_{4} \text { "); }
\]
\[
\operatorname{scanf("\% d\% d",~\& x,~\& y);~}
\]
\[
\text { printf( } \left.{ }^{\circ} \% d_{4} *_{\mathrm{L}} \%_{\mathrm{L}}==_{\mathrm{L}} \% \mathrm{~d} \backslash \mathrm{n} ", x, y, x * y\right) ;
\]
- \(\operatorname{scanf(3)}\) takes a template specifying what types of values to read and how many.
- printf(3) takes a template with holes to fill in with the values of its remaining arguments.
- \%d means scan/print an int in decimal.

\section*{Checking for input errors}
\#include <stdio.h>
```

int main()

```
```

int main()

```
\{
```

int x, y;
printf("Enter_twouintegers:ч");

```
// scanf(3) returns the number of *successful*
// conversions:
int count = scanf("\%d\%d", \&x, \&y);
if (count == 2) \{
    printf( \(\left.{ }^{\circ} \%_{\mathrm{L}} *_{\mathrm{L}} \%_{\mathrm{L}}==_{\mathrm{L}}{ }^{\circ} \mathrm{d} \backslash \mathrm{n}^{\prime}, \mathrm{x}, \mathrm{y}, \mathrm{x} * \mathrm{y}\right)\);
\} else \{
    printf("Input_error\n");
    return 1;
\}

\section*{A main function for the fib program}
src/fib_iter.c
```

\#include <stdio.h>
long fib(int n)
{ ... }
int main()
{
int input;
while (scanf("%d", \&input) == 1) {
printf("%ld\n", fib(input));
}
}

```

\section*{Structure types}

\section*{Structure types in C}

C (like BSL/ISL) uses structures to define new data types by composition of existing data types
A structure type has a name and some number of fields, each of which must be declared with a type

\section*{Syntax to define a struct type}
```

struct posn
{
double x;
double y;
};
struct circle
{
struct posn center;
double radius;
};

```

\section*{Syntax to define a struct type}
```

struct posn
{
double x;
double y;
};
struct circle
{
struct posn center;
};

```

Note that the type defined by the struct posn definition, and used for field center of struct circle is struct posn, not merely posn. (In C++ you could refer to it either way, but not in C.)

\section*{Syntax to use a structure}

Suppose we have a variable \(p\) whose type is struct posn. How do we access p's fields?

\section*{Syntax to use a structure}

Suppose we have a variable \(p\) whose type is struct posn. How do we access p's fields? p. x and p.y

\section*{Syntax to use a structure}

Suppose we have a variable \(p\) whose type is struct posn. How do we access p's fields? p.x and p.y

Let's write a function to compute the Manhattan distance between two points. Mathematically,
\[
d_{1}\left(\left(x_{1}, y_{1}\right),\left(x_{2}, y_{2}\right)\right)=\left|x_{1}-x_{2}\right|+\left|y_{1}-y_{2}\right|
\]

\section*{Syntax to use a structure}

Suppose we have a variable \(p\) whose type is struct posn. How do we access p's fields? p. \(x\) and p.y
Let's write a function to compute the Manhattan distance between two points. Mathematically,
\[
d_{1}\left(\left(x_{1}, y_{1}\right),\left(x_{2}, y_{2}\right)\right)=\left|x_{1}-x_{2}\right|+\left|y_{1}-y_{2}\right|
\]
// For the fabs(3) function:
\#include <math.h>
// Finds the Manhattan distance between two points. double manhattan_dist(struct posn \(p\), struct posn q) \{
return fabs(p.x - q.x) + fabs(p.y - q.y);
\}

\section*{Defining and initializing a structure}

Usually to get a structure in C, first you define a structure variable and then initialize it by assigning each field:
\[
\begin{aligned}
& \text { struct posn } p ; \\
& \text { p.x }=3.0 ; \\
& \text { p.y }=4.0 ; \\
& \text { struct circle c; } \\
& \text { c. center.x }=7.0 ; \\
& \text { c.center.y }=-9.2 ; \\
& \text { c. radius }=6.4 ;
\end{aligned}
\]

C won't force you to initialize all the fields, but guess what happens if you a access a field that hasn't been initialized?

\section*{Factory functions}

If you get tired of initializing structures as on the previous slide, you can always define a factory function to do the work:
```

struct circle
{
struct circle result;
result.center = center;
result.radius = radius;
return result;
}

```
make_circle(struct posn center, double radius)
(Note that functions can both take and return structure values.)

\section*{Visualizing structure value layout}
struct circle c;
c. center.x = 10.0;
c.radius = 50.0;
c.center.y = -7.0;

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struct circle c;
c.center.x = 10.0;
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c:


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struct circle c;
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c.center.y = -7.0;
c:
1.000000000 e 1

\section*{Visualizing structure value layout}
struct circle c;
c.center.x = 10.0;
c.radius = 50.0;
c.center.y = -7.0;
c:
1.000000000 e 1
5.000000000 e 1

\section*{Visualizing structure value layout}
struct circle c;
c.center.x = 10.0;
c.radius = 50.0;
c.center.y = -7.0;
c:
\(1.000000000 \mathrm{e} 1-7.000000000 \mathrm{e} 05.000000000 \mathrm{e} 1\)```


[^0]:    a fixed-width machine "integer"
    ${ }^{b}$ also a fixed-width "integer," but maybe wider

