Comp Sci 321 Programming Languages

Fall 2019

Instructor: Vincent St-Amour

Course Details

http://www.cs.northwestern.edu/~stamourv/ teaching/321-F19

(or search for "Vincent St-Amour" and follow the links)

- Slides will be posted there
- Office hours and logistic info
- Link to Piazza
- Everything, really

Course Format

- 9 homework assignments
 - Starts off gentle, but gets much harder (don't let your guard down)
 - Assignments due Fridays at 2pm
 - Grades out on Mondays (or we give you a heads up)
 - Individual submissions
- No exams
- A lot of programming, all in Racket*
 - We assume you either know Racket
 - $^{\circ}$ or can pick it up on your own
- Loosely following PLs: Application and Interpretation (PLAI)
 - $^{\circ}$ First edition
 - Link on course web page

Grading

- For each assignment, 4 possible grades:
 - **check+ (A)**: you understand this material
 - **check (B)**: you're missing some of the subtleties
 - **check- (C)**: you're missing some of the key bits
 - **0 (F)**: you should revisit that material
- Very coarse grained, and pretty strict
 - But with our resubmission policy (next slide), works well
 - $^{\rm O}$ So don't panic if you get a check- or a 0
 - $^{\circ}$ It just means you still have things to learn
- We give you some of the tests your submission failed
 - $^{\circ}$ Not all of them, though
 - $^{\circ}$ So you really want to be thorough in your own tests

Resubmission Policy

- I care that you learn the material, sooner or later
 - So you should get credit for it, even if it takes you more time
 But I still want deadlines to keep you on top of things
- If you submitted to the initial deadline, you can resubmit for up to two weeks after an assignment's deadline
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 - $^{\circ}$ That's two more chances to get feedback (and a grade)
- Resubmission grades capped halfway between "check" and "check+"
- We **expect** you to resubmit
 - Difficulty calibrated assuming you do
- I'm giving you a lot of rope; use it carefully
 Make good use of every attempt, you'll need the feedback
- Details on the course web page

Academic Integrity

- Collaboration good, plagiarism bad
 You need to understand the difference
- The work you submit must be your own
- Don't even *look* at other solutions!
 Not your colleagues'
 Not online
- We check
- We report anything suspicious to the dean

Classroom Etiquette

- Learning this (or any) material requires focus and concentration
 Let's ensure our classroom environment is conducive to that
- Laptops
 - $^{\circ}$ Laptops are fine; following along with the slides is great
 - Some laptop activities are distracting, though
 - If you're planning to do non-course-related stuff, sit in the back
 - So you don't distract your colleagues who are paying attention
- Talking
 - Asking a quick question to your neighbor is fine
 - But ask me instead, so everyone benefits from the answer
 - Continuous talking is extremely rude
 - Distracting for your colleagues around you, and for me too
 - If you want to chat, go outside

Course Staff and Office Hours



Instructor: Vincent St-Amour Wednesdays I-2 or by appointment, Mudd 3215

Peer mentors: Chloe Brown, Hakan Dingenc, Shu Han, Jenny Lam, Daniel McGrory, and Jack Wiig

See web site



TA: Spencer Florence

(Will help peer mentors at busy times.)

Why This Class? #I

Not all of you will design languages

• But all of you *might*

All these things are languages:

- Configuration file formats
- $^{\circ}$ Data formats
- Templating languages (e.g., JSX in React)
- Custom query languages (e.g., Google search queries, any "advanced search" anywhere)

They have variables, datatypes, conditionals, functions, etc. (And when they don't, often they should.)

You probably won't create the next JavaScript, but chances are you may create one of these at some point

Why This Class? #I

From a student who took the class last year:

I'm soon wrapping up my internship at [redacted], and I got to make my own compiler!

For a cloud migration, the team had to convert 100+ DAGs to Airflow. They were going to do it manually, but I volunteered and used the skills I picked up in 321 to convert Yaml files to Airflow DAGs automatically instead.

Why This Class? #I

Ok, sure, what if I do create one of these? So what?

- If you're able to recognize it as a language problem
- $^{\circ}$ You can bring to bear all we know about building languages

If you don't understand the fundamentals of languages, you'll make easily avoidable mistakes

- Strange scope rules (Python, JavaScript)
- Inconsistent handling of functions (Ruby)
- Leaky type systems (C)
- Lack of expressiveness (most configuration languages), leading to copy and paste (and bugs!)
- Etc.

In this class, you'll learn how to get the basics right

Why This Class? #2

Understanding languages makes you a more effective programmer.

- Different languages share common concepts
 - Operations on basic data (e.g., arithmetic)
 - $^{\rm O}$ Variables and scope
 - $^{\circ}$ Functions
 - $^{\circ}$ State
 - etc.
- Learn those well, and learning languages is easy!
 - New faces on familiar ideas
 - $^{\rm O}$ Languages go and come, but λ abides
- Differences between languages as variations on such concepts
 - $^{\circ}$ (Most of) the rest is (mostly) cosmetic
 - $^{\rm o}$ Then you can focus on the differences that ${\rm do}$ matter

Key Ideas of this Class

Programs are data

- ... which other programs can operate on
 - to run then (interpreters)
 - to transform them (refactoring tools)
 - to check properties about them (type checkers)
- ... which other programs can generate
 - $^{\circ}$ to make them more efficient (compilers)
 - to automate some aspects of programming (code generators)
 - $^{\circ}$ to generate infinite test cases (generative testing)
- Comes up surprisingly often in practice!

 $^{\circ}$... if you know how to look!

Key Ideas of this Class

Meta-language vs object-language

- Meta-language programs operate on programs in the object-language
- Our meta-language will be a variant of Racket: #lang plai
 Our well suited as a meta-language
- Our object-languages will be many, small, and simple
 Designed to illustrate specific concepts

This Class's Approach

Learn by building

- Both in lecture and in homeworks
- We will build interpreters
 - Interpreter = (meta-language) program that executes (object-language) programs
 - $^{\circ}$ New concept \rightarrow new object-language \rightarrow new interpreter
 - See more than one way to implement most concepts
- ... and also a few other kinds of meta-programs
 - Parsers
 - Program generators
 - Compilers
 - Type checkers

Topic Outline

- Variables and binding (substitution and deferred substitution)
- (Higher-order) functions
- Parsing (a little)
- Random testing
- Recursion
- State
- Control
- Garbage collection
- Type checking and type inference

Homework #I



On the course web page Due on Friday at 2pm

To test your prerequisites Should be very easy If not, you may not be ready

Tree traversals and manipulations will be our bread and butter So you need to master them!

Future homeworks: Also due on Fridays at 2pm

Let's dive in!

The Most Common Kinds of Program Manipulators

An *interpreter* takes a program and produces a result

• Python

 $^{\circ}$ bash

- $^{\circ}$ Racket
- ° x86 processor

Desktop calculator

• Algebra student

Good for understanding program behavior, easy to implement (our focus)

A **compiler** takes a program and produces a program

- ° gcc
- $^{\circ}$ javac
- $^{\circ}$ Racket
- ° x86 processor

Good for speed, more complex (take 322)

So, what's a **program**?

A Grammar for Algebra Programs

A grammar of Algebra in **BNF** (Backus-Naur Form):

 $\langle prog \rangle$::= $\langle defn \rangle^* \langle expr \rangle$ $\langle defn \rangle$::= $\langle id \rangle (\langle id \rangle) = \langle expr \rangle$ $\langle expr \rangle ::= (\langle expr \rangle + \langle expr \rangle)$ $| (\langle expr \rangle - \langle expr \rangle)$ $\langle id \rangle (\langle expr \rangle)$ $|\langle id \rangle$ $| \langle num \rangle$ $\langle id \rangle$::= a variable name: **f**, **x**, **y**, **z**, ... (num) ::= a number: 1, 42, 17, ...

Each **meta-variable**, such as (prog), defines a set

 $\langle id \rangle$::= a variable name: **f**, **x**, **y**, **z**, ... $\langle num \rangle$::= a number: I, 42, I7, ...

The set $\langle id \rangle$ is the set of all variable names

The set $\langle num \rangle$ is the set of all numbers

To make an example member of $\langle num \rangle$, simply pick an element from the set

 $2 \in \langle num \rangle$

 $298 \in \langle num \rangle$

$$\begin{array}{ll} \langle expr \rangle & ::= & (\langle expr \rangle + \langle expr \rangle) \\ & | & (\langle expr \rangle - \langle expr \rangle) \\ & | & \langle id \rangle (\langle expr \rangle) \\ & | & \langle id \rangle \\ & | & \langle num \rangle \end{array}$$

The set $\langle expr \rangle$ is defined in terms of other sets

We'll have to do this in steps

$$\begin{array}{ll} \langle expr \rangle & ::= & (\langle expr \rangle + \langle expr \rangle) \\ & | & (\langle expr \rangle - \langle expr \rangle) \\ & | & \langle id \rangle (\langle expr \rangle) \\ & | & \langle id \rangle \\ & | & \langle num \rangle \end{array}$$

To make an example $\langle expr \rangle$:

 $^{\rm O}$ choose one case in the grammar

 $^{\rm O}$ pick an example for each meta-variable

 $^{\circ}$ combine the examples with literal text

$$\begin{array}{l} \langle expr \rangle ::= (\langle expr \rangle + \langle expr \rangle) \\ & | (\langle expr \rangle - \langle expr \rangle) \\ & | \langle id \rangle (\langle expr \rangle) \\ & | \langle id \rangle \\ & | \langle num \rangle \end{array}$$

To make an example $\langle expr \rangle$:

 $^{\rm O}$ choose one case in the grammar

 $^{\circ}$ pick an example for each meta-variable

 $7 \in \langle num \rangle$

 $^{\rm O}$ combine the examples with literal text

 $7 \in \langle expr \rangle$

$$\begin{array}{ll} \langle expr \rangle & ::= & (\langle expr \rangle + \langle expr \rangle) \\ & | & (\langle expr \rangle - \langle expr \rangle) \\ & | & \langle id \rangle (\langle expr \rangle) \end{array}$$

To make an example $\langle expr \rangle$:

 $^{\rm O}$ choose one case in the grammar

 $^{\circ}$ pick an example for each meta-variable

 $\mathbf{f} \in \langle \mathsf{id} \rangle \qquad \qquad \mathbf{7} \in \langle \mathsf{expr} \rangle$

 $^{\circ}$ combine the examples with literal text

 $\mathbf{f}(7) \in \langle \mathsf{expr} \rangle$

$$\begin{array}{ll} \langle expr \rangle & ::= & (\langle expr \rangle + \langle expr \rangle) \\ & | & (\langle expr \rangle - \langle expr \rangle) \\ & | & \langle id \rangle (\langle expr \rangle) \end{array}$$

To make an example $\langle expr \rangle$:

 $^{\rm O}$ choose one case in the grammar

 $^{\circ}$ pick an example for each meta-variable

 $\mathbf{f} \in \langle \mathsf{id} \rangle \qquad \qquad \mathbf{f}(7) \in \langle \mathsf{expr} \rangle$

 $^{\circ}$ combine the examples with literal text

 $\mathbf{f}(\mathbf{f}(7)) \in \langle \mathsf{expr} \rangle$

 $\langle prog \rangle$::= $\langle defn \rangle^* \langle expr \rangle$ $\langle defn \rangle$::= $\langle id \rangle (\langle id \rangle) = \langle expr \rangle$ $\mathbf{f}(\mathbf{x}) = (\mathbf{x} + 1) \in \langle defn \rangle$

* is a Kleene star: means 0 or more

To make a (prog) pick some number of (defn)s

 $(\mathbf{x} + \mathbf{y}) \in \langle \operatorname{prog} \rangle$

 $f(\mathbf{x}) = (\mathbf{x} + 1)$ $g(\mathbf{y}) = f((\mathbf{y} - 2)) \in \langle \operatorname{prog} \rangle$ g(7)

So, what's a language, then?

A programming language is defined by

- a grammar that describes what programs are possible
- rules for evaluating any such program to produce a result

For example, algebra evaluation is defined in terms of evaluation steps:

 $(2 + (7 - 4)) \rightarrow (2 + 3) \rightarrow 5$

So, what's a language, then?

A programming language is defined by

- a grammar that describes what programs are possible
- rules for evaluating any such program to produce a result

For example, algebra evaluation is defined in terms of evaluation steps:

 $f(\mathbf{x}) = (\mathbf{x} + \mathbf{I})$ $f(\mathbf{I}0) \longrightarrow (\mathbf{I}0 + \mathbf{I}) \longrightarrow \mathbf{I}\mathbf{I}$

Evaluation

 Evaluation (→) is defined by a set of pattern-matching rules:

$$(2 + (7 - 4)) \rightarrow (2 + 3)$$

due to the rule
... (7 - 4) ... \rightarrow ... 3 ...

Evaluation

 Evaluation (→) is defined by a set of pattern-matching rules:

> $f(\mathbf{x}) = (\mathbf{x} + \mathbf{I})$ $f(\mathbf{I}0) \longrightarrow (\mathbf{I}0 + \mathbf{I})$

> > due to the rule

$$\cdots \langle id \rangle_{I} (\langle id \rangle_{2}) = \langle expr \rangle_{I} \cdots$$

$$\cdots \langle id \rangle_{I} (\langle expr \rangle_{2}) \cdots \rightarrow \cdots \langle expr \rangle_{3} \cdots$$

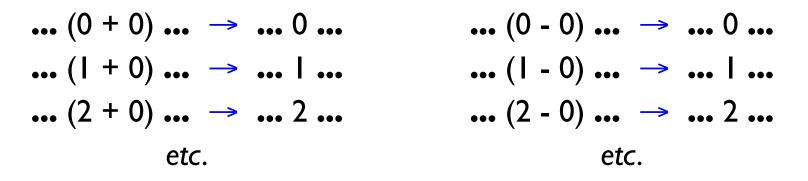
where $\langle expr \rangle_3$ is $\langle expr \rangle_1$ with $\langle id \rangle_2$ replaced by $\langle expr \rangle_2$

Rules for Evaluation

• Rule I: one pattern

 $... \langle id \rangle_{I} (\langle id \rangle_{2}) = \langle expr \rangle_{I} ...$ $... \langle id \rangle_{I} (\langle expr \rangle_{2}) ... \rightarrow ... \langle expr \rangle_{3} ...$ $where \langle expr \rangle_{3} is \langle expr \rangle_{I} with \langle id \rangle_{2} replaced by \langle expr \rangle_{2}$

• **Rules 2 -** ∞ : special cases



When the interpreter is a program instead of an Algebra student, the rules look a little different

Action Items

- Sign up for Piazza
- Brush up your Racket
- Read the docs for the PLAI language (comes with Racket) http://docs.racket-lang.org/plai/plai-scheme.html
- · Do Homework I