

Higher-Order Functions (Part I)

Why Functions as Values

- Abstraction is easier with functions as values
 - abstract over **add** and **sub** cases
 - **filter**, **map**, etc.
- What are objects? Callbacks?
- Separate **defun** form becomes unnecessary
 - ```
{defun {f x} {+ 1 x}}
{f 10}
```

  
⇒  

```
{with {f {fun {x} {+ 1 x}}}
 {f 10}}
```

# FWAE Grammar, Almost

```
<FWAE> ::= <num>
| {+ <FWAE> <FWAE>}
| {- <FWAE> <FWAE>}
| {with {<id> <FWAE>} <FWAE>}
| <id>
| {<id> <FWAE>}
| {fun {<id>} <FWAE>}
```

?

NEW

# FWAE Evaluation

10  $\Rightarrow$  10

{+ 1 2}  $\Rightarrow$  3

{- 1 2}  $\Rightarrow$  -1

{with {x 7} {+ x 2}}  $\Rightarrow$  {+ 7 2}  $\Rightarrow$  9

y  $\Rightarrow$  *free identifier*

{fun {x} {+ 1 x}}  $\Rightarrow$

{fun {x} {+ 1 x}}

Result is not always a number!

; interp : FWAE? ...  $\rightarrow$  FWAE-Value?

# FWAE Evaluation

Let's think in terms of substitution, for simplicity

```
 {with {y 10} {fun {x} {+ y x}}}
⇒ {fun {x} {+ 10 x}}
```

```
{with {f {fun {x} {+ 1 x}}}
 {f 3}}
⇒ {{fun {x} {+ 1 x}} 3}
```

Doesn't match the grammar for <FWAE>

# FWAE Grammar

```
<FWAE> ::= <num>
| {+ <FWAE> <FWAE>}
| {- <FWAE> <FWAE>}
| {with {<id> <FWAE>} <FWAE>}
| <id>
| {<id> <FWAE>}
| {fun {<id>} <FWAE>}
| {<FWAE> <FWAE>}
```

NEW

NEW

# FWAE Evaluation

`{with {f {fun {x} {+ 1 x}}} {f 3}}`

$\Rightarrow$  `{{fun {x} {+ 1 x}} 3}`

$\Rightarrow$  `{+ 1 3}`  $\Rightarrow$  `4`

`{{fun {x} {+ 1 x}} 3}`  $\Rightarrow$  `{+ 1 3}`  $\Rightarrow$  `4`

`{1 2}`  $\Rightarrow$  *expected function*

`{+ 1 {fun {x} 10}}`  $\Rightarrow$  *expected number*

# FWAE Datatype

```
(define-type FWAE
 [num (n number?)]
 [add (lhs FWAE?)
 (rhs FWAE?)]
 [sub (lhs FWAE?)
 (rhs FWAE?)]
 [with (name symbol?)
 (named-expr FWAE?)
 (body FWAE?)]
 [id (name symbol?)]
 [fun (param-name symbol?)
 (body FWAE?)]
 [app (fun-expr FWAE?)
 (arg-expr FWAE?)])

(test (parse '{fun {x} {+ x 1}})
 (fun 'x (add (id 'x) (num 1))))
```



# FWAE Datatype

```
(define-type FWAE
 [num (n number?)]
 [add (lhs FWAE?)
 (rhs FWAE?)]
 [sub (lhs FWAE?)
 (rhs FWAE?)]
 [with (name symbol?)
 (named-expr FWAE?)
 (body FWAE?)]
 [id (name symbol?)]
 [fun (param-name symbol?)
 (body FWAE?)]
 [app (fun-expr FWAE?)
 (arg-expr FWAE?)])
```

```
(test (parse '{{fun {x} {+ x 1}} 10})
 (app (fun 'x (add (id 'x) (num 1))) (num 10)))
```

# FWAE-Value

```
(define-type FWAE-Value
 [numV (n number?)]
 [funV (param-name symbol?)
 (body FWAE?)])
```

# FWAE Interpreter

```
; interp : FWAE? -> FWAE-Value?
(define (interp an-fwae)
 (type-case FWAE an-fwae
 [num (n) (numV n)]
 [add (l r) (num+ (interp l) (interp r))]
 [sub (l r) (num- (interp l) (interp r))]
 [with (name named-expr body)
 (interp (subst body
 name
 (interp named-expr)))]
 [id (name) (error 'interp "free identifier")]
 [fun (param-name body)
 (funV param-name body)]
 [app (fun-expr arg-expr)
 (define fun-val (interp fun-expr))
 (interp (subst (funV-body fun-val)
 (funV-param-name fun-val)
 (interp arg-expr)))]))
```

# Add and Subtract

```
; num+ : FWAE-Value? FWAE-Value? -> FWAE-Value?
(define (num+ x y)
 (numV (+ (numV-n x) (numV-n y))))
; num- : FWAE-Value? FWAE-Value? -> FWAE-Value?
(define (num- x y)
 (numV (- (numV-n x) (numV-n y))))
```

Better:

```
; num-op :
; (number? number? -> number?) ->
; (FWAE-Value? FWAE-Value? -> FWAE-Value?)
(define (num-op op)
 (lambda (x y)
 (numV (op (numV-n x) (numV-n y)))))

(define num+ (num-op +))
(define num- (num-op -))
```

## FWAE Subst

```
; subst : FWAE? symbol? FWAE-Value? -> FWAE?
(define (subst exp sub-id val)
 (type-case FWAE exp
 ...
 [id (name)
 (cond
 [(equal? name sub-id)
 (type-case FWAE-Value val
 [numV (n) (num n)]
 [funV (param-name body)
 (fun param-name body)]]
 [else exp])])
 ...))
```

# FWAE Subst

```
; subst : FWAE? symbol? FWAE-Value? -> FWAE?
(define (subst exp sub-id val)
 (type-case FWAE exp
 ...
 [app (f arg)
 (app (subst f sub-id val)
 (subst arg sub-id val))]
 [fun (param-name body)
 (if (equal? sub-id param-name)
 exp
 (fun param-name
 (subst body sub-id val)))]))
```

# FWAE Subst

Beware: with the implementation on the previous slide,

```
(subst {with {y 10} z}
 'z
 {fun {x} {+ x y}})
⇒ {with {y 10} {fun {x} {+ x y}}}
```

- That **y** in the function used to be a free identifier
- and now it's bound! That can't be right...
- We'll see how to fix this

# No More With

Compare the **with** and **app** implementations:

```
(define (interp an-fwae)
 (type-case FWAE an-fwae
 ...
 [with (name named-expr body)
 (interp (subst body
 name
 (interp named-expr)))]
 ...
 [app (fun-expr arg-expr)
 (define fun-val (interp fun-expr))
 (interp (subst (funV-body fun-val)
 (funV-param-name fun-val)
 (interp arg-expr)))]))
```

The **app** case does everything that **with** does



# No More With

```
{with {x 10} x}
```

is the same as

```
{{fun {x} x} 10}
```

In general,

```
{with {<id> <FWAE>1} <FWAE>2}
```

is the same as

```
{{fun {<id>} <FWAE>2} <FWAE>1}
```

Aside: IIFEs in JavaScript

So let's just get rid of `with`

Don't worry, we'll bring it back

# FAE Grammar

```
<FAE> ::= <num>
 | {+ <FAE> <FAE>}
 | {- <FAE> <FAE>}
 | {with {<id> <FAE>} <FAE>}
 | <id>
 | {fun {<id>} <FAE>}
 | {<FAE> <FAE>}
```

# F AE Interpreter

```
; interp : FAE? -> FAE-Value?
(define (interp a-fae)
 (type-case FAE a-fae
 [num (n) (numV n)]
 [add (l r) (num+ (interp l) (interp r))]
 [sub (l r) (num- (interp l) (interp r))]
 [id (name) (error 'interp "free identifier")]
 [fun (param-name body) (funV param-name body)]
 [app (fun-expr arg-expr)
 (define fun-val (interp fun-expr))
 (interp (subst (funV-body fun-val)
 (funV-param-name fun-val)
 (interp arg-expr))))]))
```

# Where to?

- This FAE language will be our "base camp" for most of the rest of the quarter
  - We'll use it to demonstrate concepts
  - We'll extend it in various ways
  
- Some loose ends we need to tie up
  - Fixing substitution
  - Bringing **with** back
  - Moving to deferred substitution
  - Bringing recursion back (we "lost" it along the way)