Copying Garbage Collection

Two-Space Copying Collectors

A **two-space** copying collector compacts memory as it collects, making allocation easier.

Allocator:

- Partitions memory into to-space and from-space
- Allocates only in **to-space** (in order, c.f. non-collecting)

Collector:

- Starts by swapping **to-space** and **from-space**
- Coloring gray ⇒ copy from *from-space* to *to-space*
- Choosing gray records ⇒ go through the new to-space, update pointers



Left = from-space Right = to-space























- Cool diagrams, bro
- But what does that look like for an actual heap?
- Like, say, in plai/gc2?
- So let's go through a more concrete example
- But the actual plai/gc2 implementation is your job for HW8

The Setup

- Each object in memory starts with a tag
 Just like in plai/gc2
- Tags tell us how to interpret the heap cells that follow
 - $^{\circ}$ How many cells are part of the object?
 - Which cells hold pointers?
 - Which cells hold flat data?
 - O Just like in plai/gc2

The Setup

- The kinds of objects we'll be dealing with are simplified variants of the ones in plai/gc2
- Flat data will be integers only, to keep things simple
- Tag i: one integer

Simpler variant of 'flat

• Tag b: one pointer

• Simpler variant of 'cons (like a box)

- Tag c: one integer, then one pointer
 Simpler variant of 'clos
- Tag **f**: forwarding pointer (one pointer)

The Strategy

- Traverse the heap, starting at the roots, using breadth-first search $^{\circ}$ In contrast, mark-and-sweep uses depth-first
- Visiting a node = marking it gray
 - \circ = copying from the from-space to the to-space
 - \circ + leaving a forwarding pointer behind in the from-space

The Strategy

- Maintain a queue of the gray nodes in the to-space
 - $^{\circ}$ Marking a node gray \rightarrow adding it to the queue
 - $^{\circ}$ Taking a node out of the queue \rightarrow marking it black
- Use that queue to keep track of the BFS

• Invariant:

- ° objects in the queue have pointers to the from-space;
- $^{\circ}$ objects outside the queue (black) have pointers to the to-space
- Represent the queue as two pointers into the to-space
 - $^{\rm O}$ Increment the end pointer when enqueuing
 - $^{\rm O}$ Increment the front pointer when dequeuing
 - When the two pointers come together, queue is empty
 - $^{\circ}$ l.e., we're done

- 26-cell memory (13 cells per space), 2 roots
 - Tag i: one integer
 - Tag b: one pointer
 - $^{\rm O}$ Tag c: one integer, then one pointer



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Two-Space Pros and Cons

- Doesn't suffer from fragmentation
- Time cost proportional to live data (not garbage!)
- Allocation is simple, just bump a pointer
- Collection doesn't require much state (handful of pointers, no stack)

- Only half the heap is in use at any time

 Not a big deal when combined with generational collection
 - Still "stop the world"

Tips for Debugging Homework 8

You may need to do a lot of debugging, and it may be painful.

- Write your heap checker first.
- Make the heap smaller to trigger GC more often.
- To stress-test your GC when debugging, GC on every allocation (not just when you run out of space).
- Pause to look at the heap when necessary (i.e., call read).
- Make sure you're not forgetting any roots.

Further reading

- GC first appeared circa 1958 (original LISP)
- Went mainstream with Java in the 90s
- Tremedous amount of work: new techniques, improvements, etc.
- Still an active research area to this day

Good reference: Uniprocessor Garbage Collection Techniques, by Wilson ftp://ftp.cs.utexas.edu/pub/garbage/gcsurvey.ps